

# Good LDPC decoders with exchanged messages of size 1, 2 and 3 bits.

Workshop on Coding for Future Optical Communications  
(Munich, 27 February 2018)

Emmanuel Boutillon

in collaboration with Chris Winstead, David Declercq and Franklin Cochachin

07/12/2018



1

# Outline

- Syndrome Bit Flipping:  
BSC channel, no extrinsic, (Bit Flipping Algorithm).
- Noisy Gradient Bit Flipping Decoder:  
AWGN Channel, no extrinsic (Bit Flipping Algorithm).
- Sign-Preserving algorithm.  
AWGN Channel, message passing algorithm

# Outline

- **Syndrome Bit Flipping:**  
BSC channel, no extrinsic, (Bit Flipping Algorithm).
- Noisy Gradient Bit Flipping Decoder:  
AWGN Channel, no extrinsic (Bit Flipping Algorithm).
- Sign-Preserving algorithm.  
AWGN Channel, message passing algorithm



# Syndrome Bit Flipping

For LDPC code IEEE802.3 (size 2048, rate 0.84,  $d_v = 6$ ,  $d_c = 32$ ):

20 errors among 2048 bits generate in average 60 unsatisfied parities among 384.

$$\binom{2048}{20} = 10^{48} \quad \text{but} \quad \binom{384}{60} = 10^{70}$$

Syndrome can be viewed as a compressed representation of the error pattern:

It should be possible to consider **only the syndrome** to suppress the error floor.

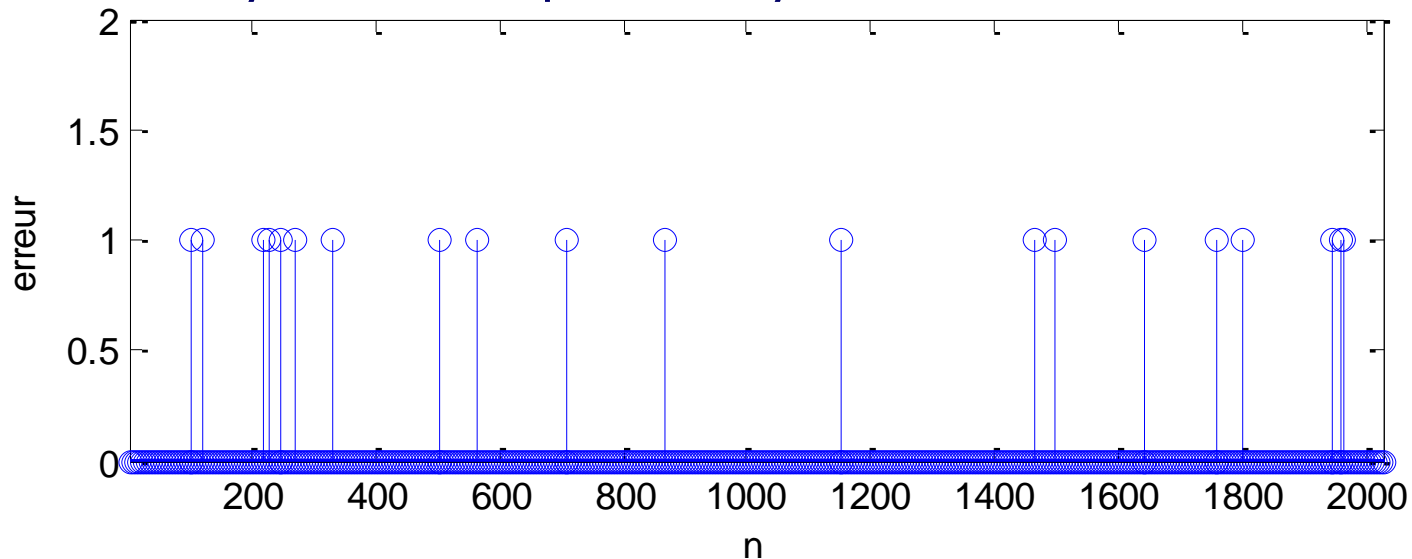
Old idea [1]: used the number of unsatisfied check connected to a variable to flip its state.

[1] Raul Benet, Adriaan De Lind Van Wijngaarden, Ralf Dohmen, Thomas Richardson, Rudiger Urbanke, "Iterative decoding of low-density parity-check (LDPC) codes", Patent EP1643653A1, 5 april 2006.

# Example of SBF on a real case

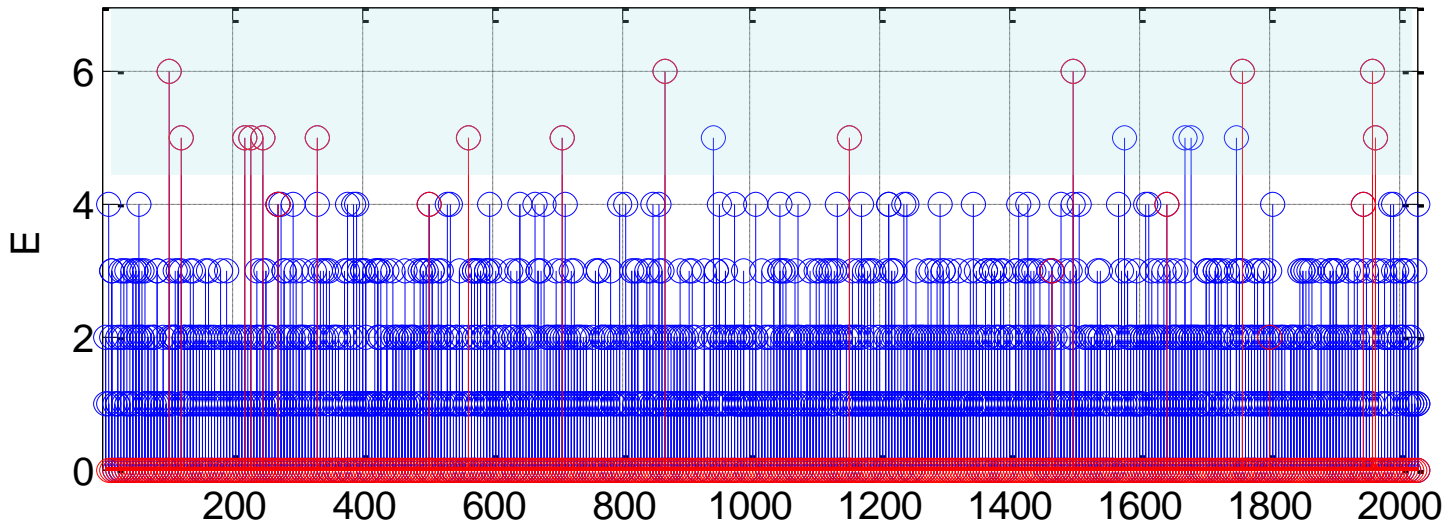
First decoder failed: the final state  $x$  doesn't respect  $H.x = 0$ .

$x = c + \text{err} \Rightarrow S = H.x \text{ mod } 2 = H(c+\text{err}) \text{ mod } 2 = 0 + H.\text{err} \text{ mod } 2$   
 $\Rightarrow$  syndrom  $S$  depends only of the error vector  $\text{err}$ .



20 errors, 92 checks non fulfilled ( $\text{sum}(S) = 92$ )

# Value of E at the first iteration



Flip bits with  $E > 4 \Rightarrow 19$  bits flipped :  
 14 that was in error  
 5 that was correct

Only  $20 - 14 + 5 = 11$  errors after the first iteration.

Sequence  $\theta =$

		4	4	4	3	3	
nb_err	=	20	11	6	5	3	0
nb_synd	=	92	44	22	18	14	0



# Impact of threshold sequence

Threshold sequence [4 4 4 3 3 ] =>

nb_err	=	20	11	6	5	3	0
nb_synd	=	92	44	22	18	14	0

Threshold sequence [4 4 4 4 4 3 3]) =>

nb_err	=	20	11	6	5	5	5	3	0
nb_synd	=	92	44	22	18	18	18	14	0

Threshold sequence [3 3 3 3 ] =>

nb_err	=	20	54	157	542	941
nb_synd	=	92	118	172	206	202

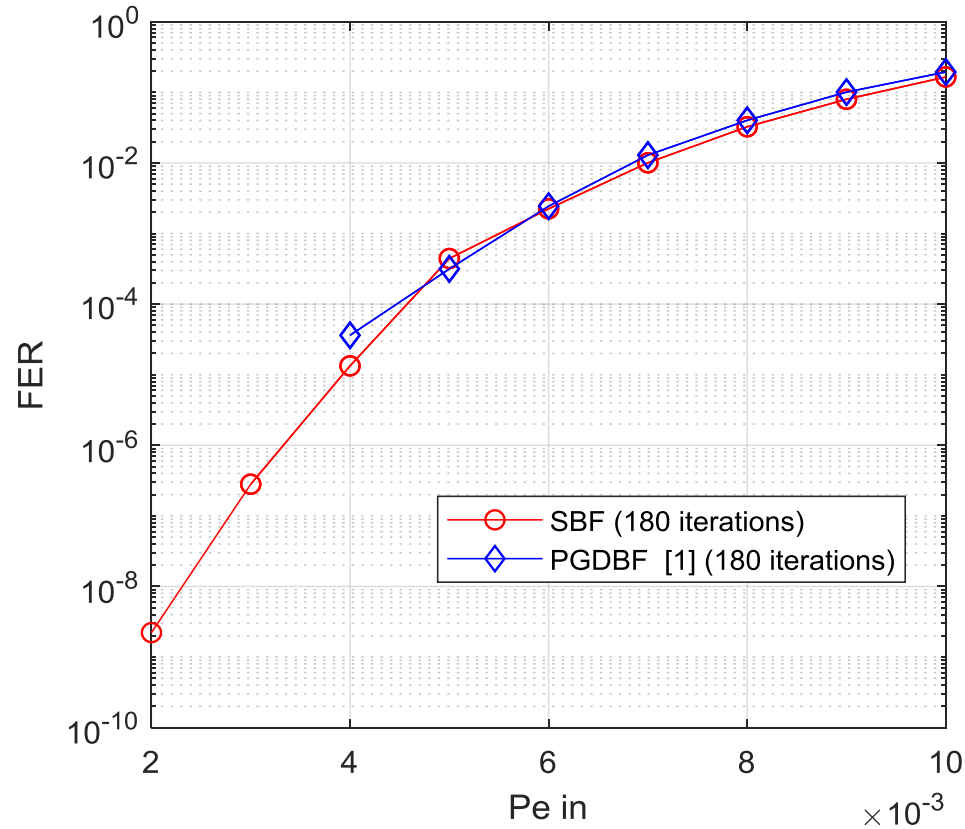
NOTE: there is error patterns solved by [3 3 3 3 ] but not by [4 4 4 3 3] !

A threshold sequence is « a decoding key ».

Each key can resolve a subset of error patterns.

# SBF alone in BSC for IEEE 802.3

$(d_v, d_c) = (6, 32), N = 2048$



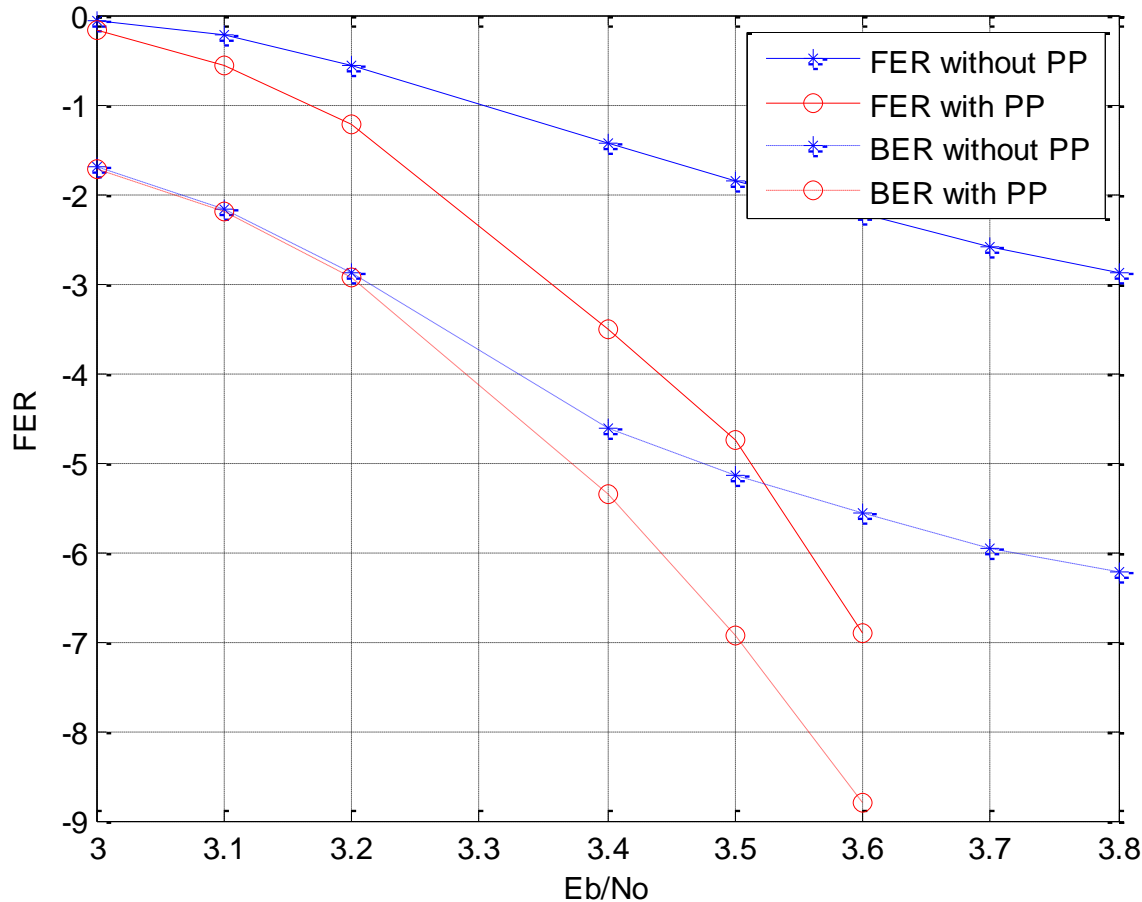
SBF alone can outperform PGDBF at low FER

[1] K. Le et al. "A novel high-throughput, low-complexity bit-flipping decoder for LDPC codes,"

# Example of SBF Post-Processing

$d_v = 5, d_c = 20, N = 10240$

Code H512, with and without post-processing (PP)



Initial algorithm that converge fast toward low weight error pattern.

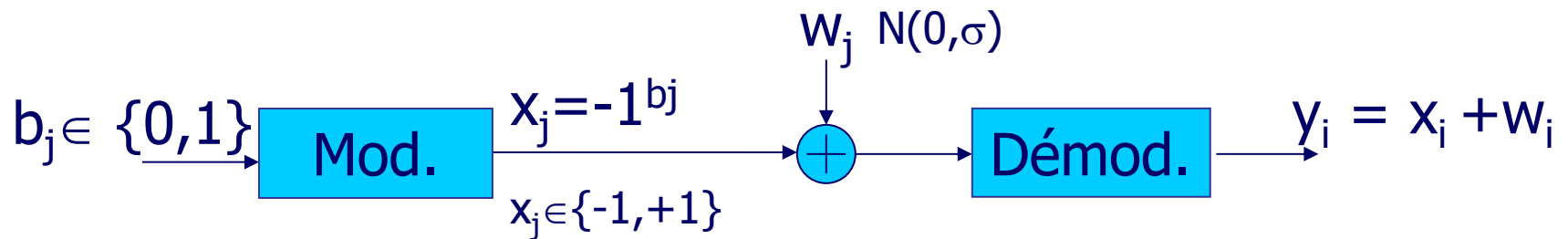
Then post-processing with SBF to clear remaining error Patern.  
(only 11 iterations of SBF in this example).

# Outline

- Syndrome Bit Flipping:  
BSC channel, no extrinsic, (Bit Flipping Algorithm).
- Noisy Gradient Bit Flipping Decoder:  
AWGN Channel, no extrinsic (Bit Flipping Algorithm).
- Sign-Preserving algorithm.  
AWGN Channel, message passing algorithm

# Optimum decoding

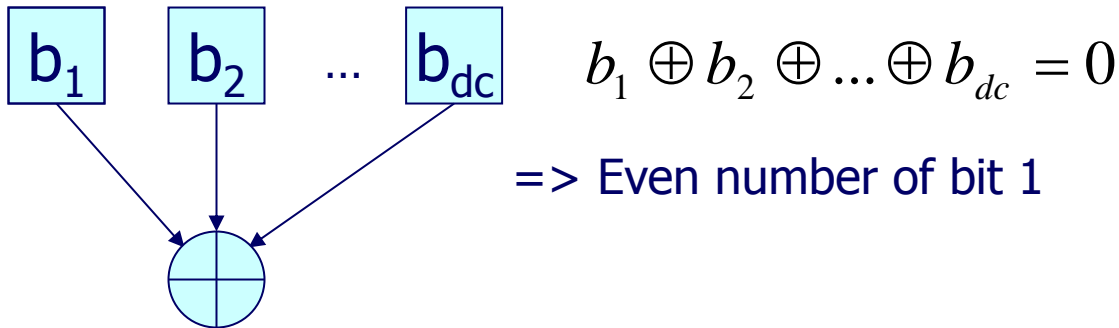
Bits of a codeword are transmitted in a noisy channel.



Decision with channel value:  $x_i = \text{sgn}(y_i)$

ML decoder (optimal decoder):  $\hat{x} = \arg \max_{x \in C} \left\{ \sum_{i=1}^n x_i y_i \right\}$

# Check with sign (BPSK 0=>1, 1=>-1)



If there is no error, then, for all  $j = 1 \dots m$   $S_j = \prod_{k \in N(j)} x_k = +1$

$$\hat{x} = \arg \max_{x \in C} \left\{ \sum_{i=1}^n x_i y_i \right\} \Leftrightarrow \hat{x} = \arg \max_{x \in C} \left\{ \sum_{i=1}^n x_i y_i + \sum_{j=1}^m S_j \right\}$$

Try to maximise the energy function  $E(x_1, \dots, x_n) = \sum_{i=1}^n x_i y_i + \sum_{j=1}^m S_j$



# Gradient bit flipping algorithm

Let assume that  $x$  is not a codeword

Question: do we increase energy function flipping bit  $x_l$ ?

If  $x_l$  flips to  $\bar{x}_l$  then

$$E_{old} \Rightarrow E_{new} = E_{old} + (\bar{x}_l - x_l) \frac{dE}{dx_l} = E_{old} - 2x_l \frac{dE}{dx_l}$$

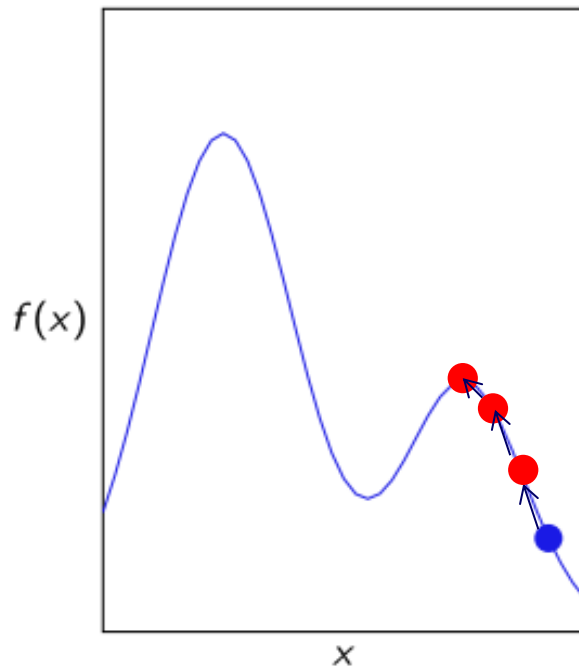
Thus, energy increases only if  $E_l = x_l \frac{dE}{dx_l} < 0$  .

with  $E_l = x_l \frac{dE}{dx_l} = x_l y_l + \sum_{j \in N(l)} S_j$

Parallel update rule: flip the bits / verifying that  $E_j < \theta$ , with  $\theta$  a negative threshold.

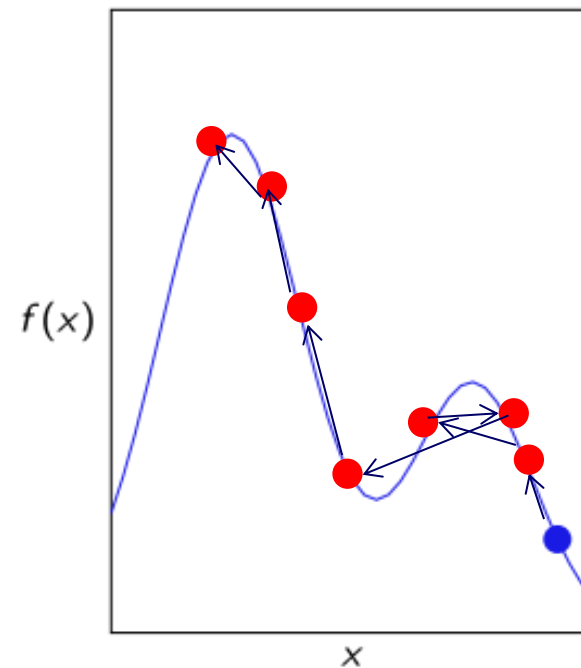


# Problem with Gradient bit flipping algorithm



Problem: algorithm sticks in a local maxima.

$$E_i = x_i y_i + \sum_{j \in M(i)}^m S_j$$



Solution to escape the local maxima: add perturbation.

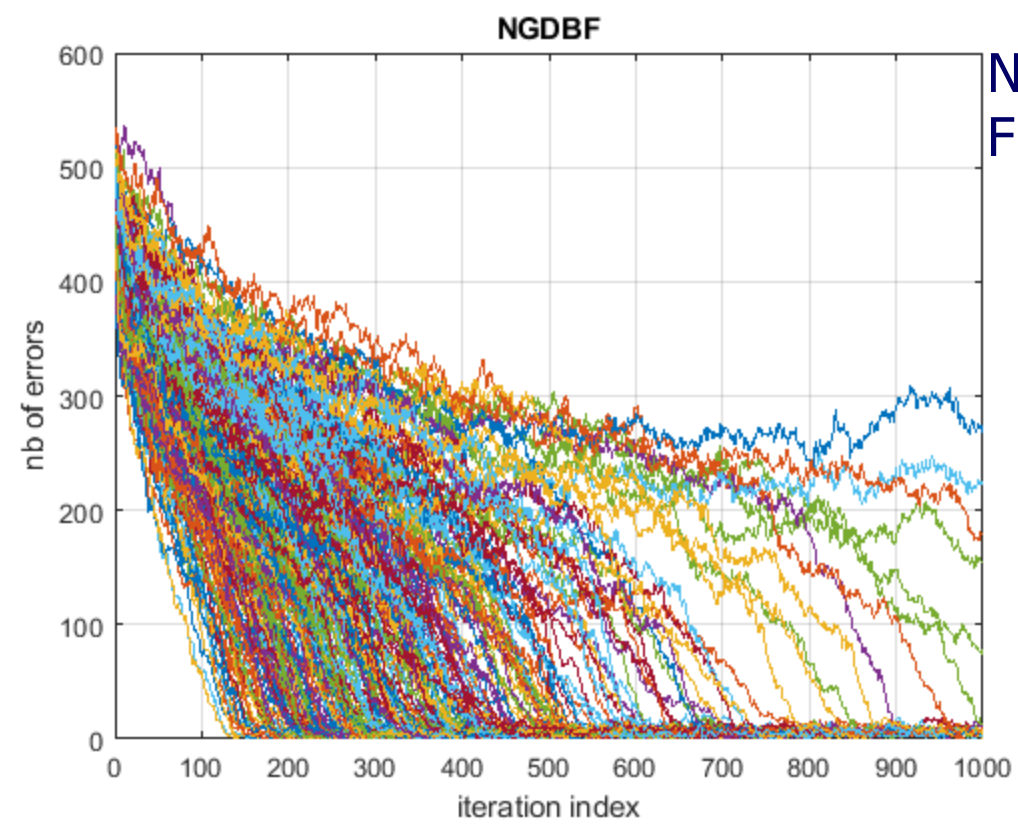
$$E_i = x_i y_i + \sum_{j \in M(i)}^m S_j + q_i$$

AWGN



# NGDBF alone

( $d_v=5, d_c = 20$ ),  
 $N = 10240$ ,  
 $E_b/N_0 = 3.2$  dB



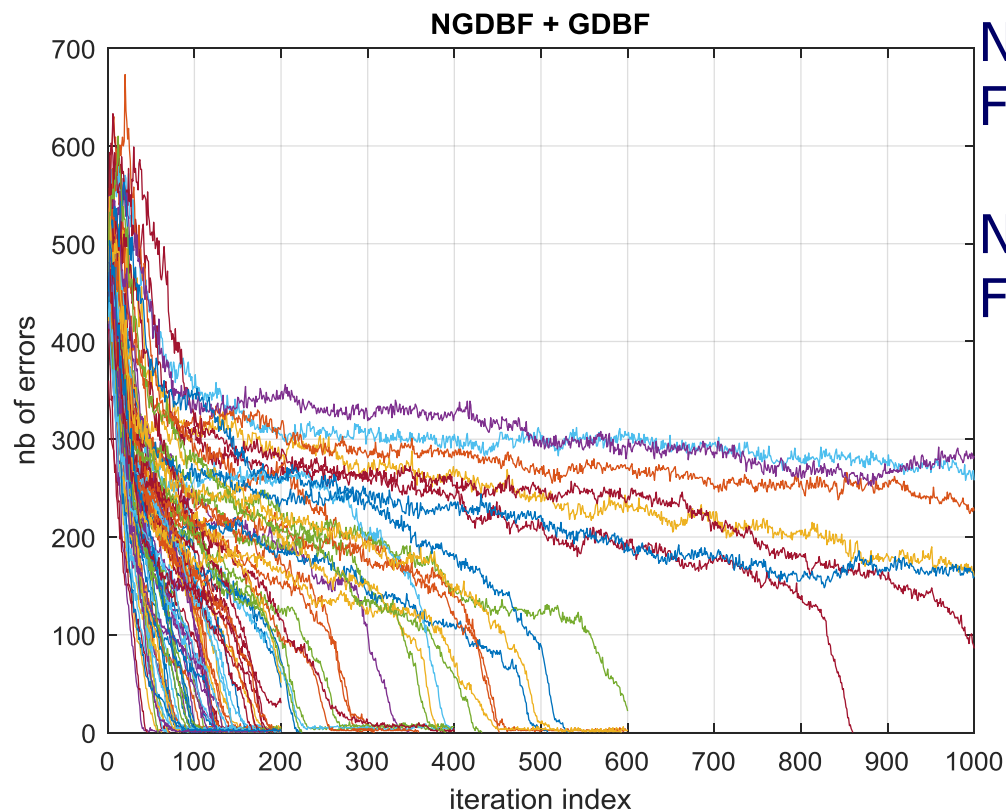
NGDBF + SBF  
FER = 0.0278

Performance reference



# Mixte NGDBF and GDBF

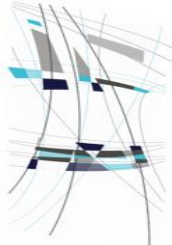
( $d_v=5, d_c = 20$ ),  
 $N = 10240$ ,  
 $E_b/N_0 = 3.2$  dB



NGDBF + SBF  
FER = 0.0278

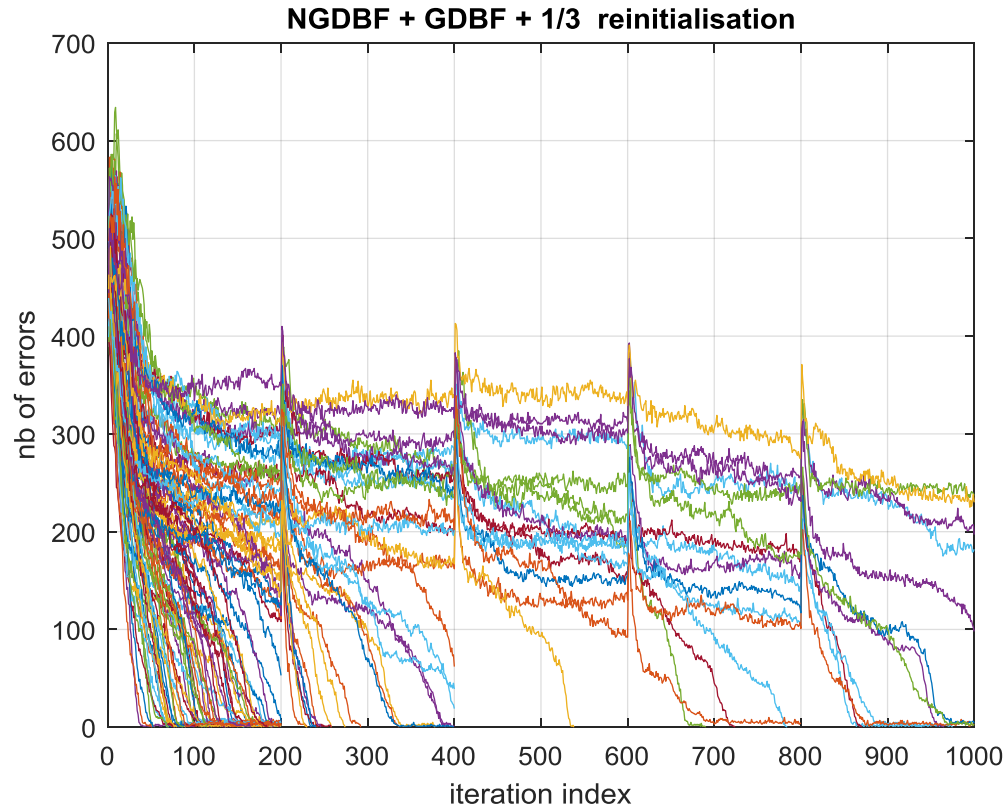
NGDBF/GDBF + SBF  
FER = 0.0585

Faster convergence, but FER degradation



# Every 200 cycles, 1/3 of states take their initial value.

( $d_v=5, d_c = 20$ ),  
 $N = 10240$ ,  
 $E_b/N_0 = 3.2$  dB

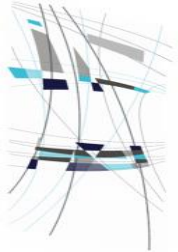


NGDBF  
FER = 0.0278

NGDBF + GDBF  
FER = 0.0585

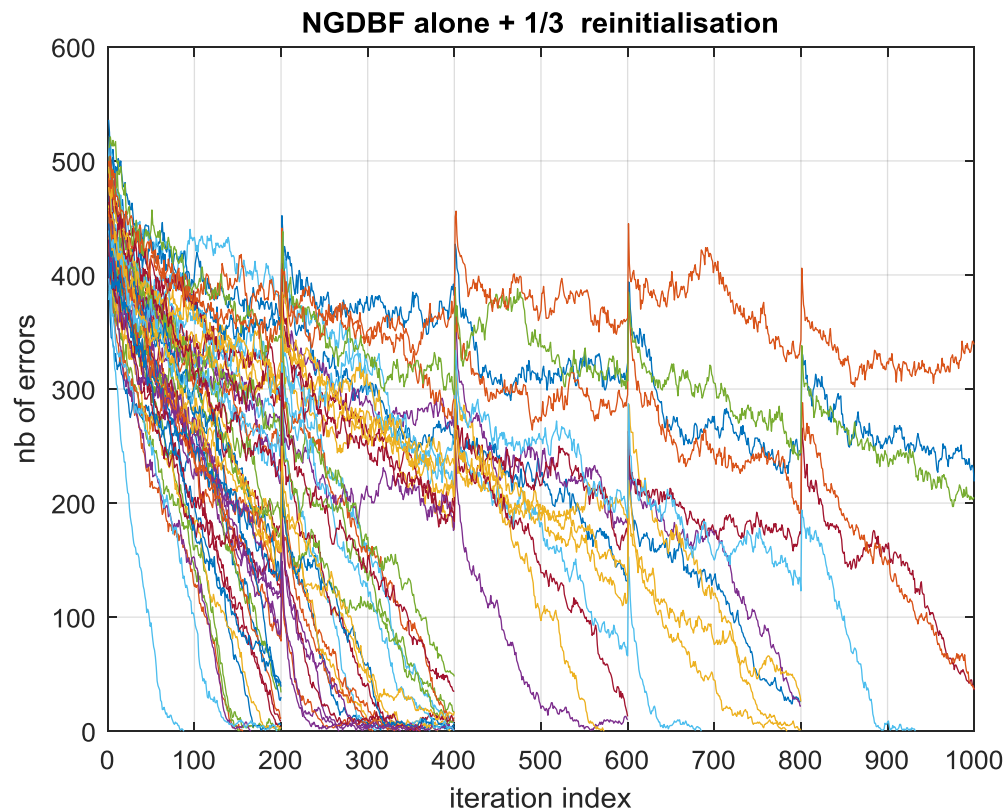
NGDBF + GDBF  
+ partial restart  
FER = 0.0188

Faster convergence, better performance.



# Réinitialisation partielle + NGDBF simple.

( $d_v=5, d_c = 20$ ),  
 $N = 10240$ ,  
 $E_b/N_0 = 3.2$  dB



NGDBF  
FER = 0.0278

NGDBF + GDBF  
FER = 0.0585

NGDBF + GDBF  
+ partial restart  
FER = 0.0188

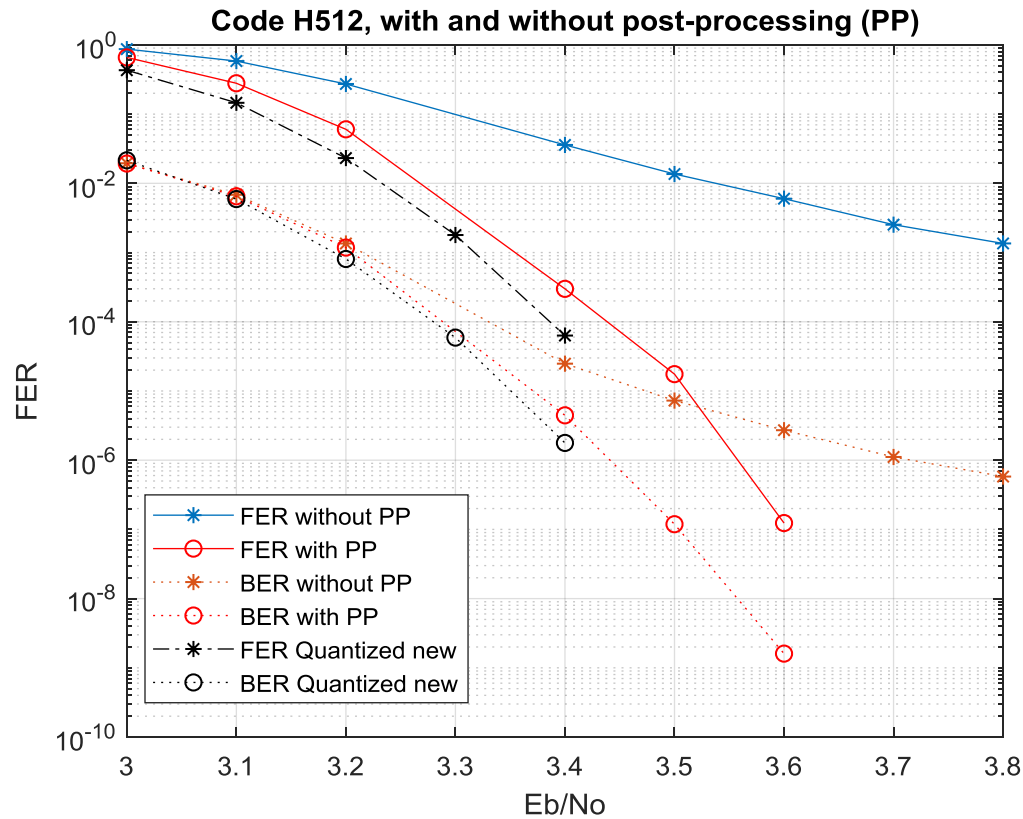
NGDBF alone  
+ Partial restart  
FER = 0.0333

Lower performance because of lower convergence



# Nouveau NGDBF

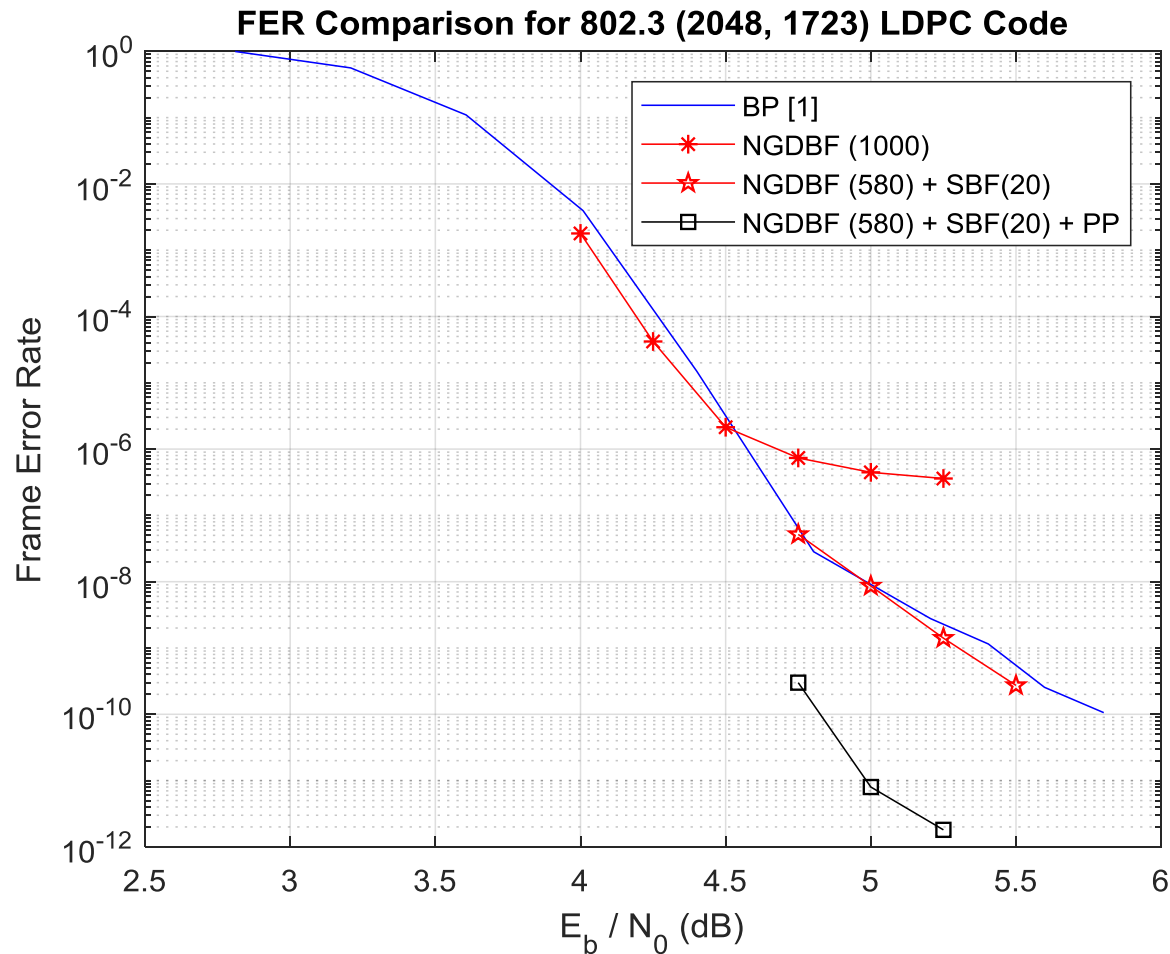
$d_v = 5, d_c = 20, N = 10240$



1000 iterations  
+ SBF

# NGDBF + SBF for IEEE 802.3

$(d_v, d_c) = (6, 32), N = 2048$





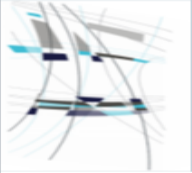
# Outline

- Syndrome Bit Flipping:  
BSC channel, no extrinsic, (Bit Flipping Algorithm).
- Noisy Gradient Bit Flipping Decoder:  
AWGN Channel, no extrinsic (Bit Flipping Algorithm).
- Sign-Preserving algorithm.  
AWGN Channel, message passing algorithm



# Outline

- Syndrome Bit Flipping:  
BSC channel, no extrinsic, (Bit Flipping Algorithm).
- Noisy Gradient Bit Flipping Decoder:  
AWGN Channel, no extrinsic (Bit Flipping Algorithm).
- Sign-Preserving algorithm.  
AWGN Channel, message passing algorithm



# Oups!

Let's us shift toward  
slides generated by latex...